

## Solutions to Homework 3

3.2

hex	binary
0x5ED4	0101 1110 1101 0100
-0x07A4	-0000 0111 1010 0100
0x5730	0101 0111 0011 0000

Note that these are in sign-magnitude format, so you shouldn't be converting to two's complement.

3.4

octal	binary
4365	100 011 110 101
-3412	-011 100 001 010
0753	000 111 101 011

3.7

185 = 10111001. But this is in sign-magnitude, so that really means = -0111001 in 7 bits.  
 122 = 01111010 = +1111010 in 7 bits. What we have, then, is actually a 7-bit subtraction:

$$\begin{array}{r} 1111010 \\ -0111001 \\ \hline 1000001 = 65_{10} \end{array}$$

This fits into 8 bits: 01000001. No overflow.

3.20

Two's complement: 201,326,592

Unsigned: Same as above

3.21

000011	000000000000000000000000
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This turns out to be the instruction: `jal 0`

3.22

0	00011000	000000000000000000000000
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Use formula:  $(-1)^s \times (1 + M) \times 2^{exp-bias} = 1.0 \times 2^{-103}$

**BUT**, this is **not** in decimal! This is in base 2!! Remember that the formula just plugs in the mantissa, which is in base 2, plus we're multiplying by 2, which *is* the base. Thus, multiply it out and the **decimal** answer is  $9.860761315 \times 10^{-32}$ .

As an example, suppose an answer came out to  $1.1 \times 2^1$ . This is really  $1.1_2 \times 2 = 1.5_{10} \times 2 = 3_{10}$ , not  $2.2_{10}$ .

3.23

$$63.25 = 111111.01 = 1.1111101 \times 2^{5+127=132=10000100}$$

0	10000100	111110100000000000000000
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3.29

Oh, boy.

$$2.6125 \times 10^1 = 26.125 = 11010.001_2$$

$$4.150390625 \times 10^{-1} = 0.4150390625 = 0.0110101001_2$$

Now normalize the base 2 values:

$$1.1010001 \times 2^4$$

$$1.10101001 \times 2^{-2}$$

To add, align the decimals by right-shifting the smaller number to get the same exponents. This means we have to right-shift the bottom value by 6. This gives us:

$$\begin{array}{r} 1.1010001000\ 0000 \times 2^4 \\ +\ 0.0000011010\ 1001 \times 2^4 \\ \hline 1.1010100010\ 1001 \times 2^4 \end{array}$$

But we have a machine with 16 bits, so the mantissa is only 10 bits. That means we have the extra bit pattern 1001 hanging off the right side. Since we have a guard, round, and sticky bit, that means we can only handle three of those bits while adding. Because there's one bit to the right, the sticky bit becomes 1, and we have:

$$\begin{array}{r} 1.1010001000\ 000 \times 2^4 \\ +\ 0.0000011010\ 101 \times 2^4 \\ \hline 1.1010100010\ 101 \times 2^4 \end{array}$$

Finally, to store the final number, the pattern 101 indicates **round up**, so the last of the 10 bits we can store will change to 1:

$$1.1010100011 \times 2^4$$

This works out to 26.546875.

Additional problems:

1.  $9FFE\ FFB_{16} = 1001\ 1111\ 1111\ 1110\ 1111\ 1111\ 1011\ 0010 = 2,684,288,946_{10}$
2.  $1111\ 1111\ 1111\ 1111\ 1111\ 1110\ 1111\ 1110 = -258_{10}$
3.  $0000\ 0000\ 0001\ 1111\ 0111\ 1111\ 1111\ 0111 = 2,064,375_{10}$

4.

M	A	Q	Q <sub>-1</sub>	Count	Operation
10011	00000	01000	0	5	initial
	00000	00100	0	4	arithmetic shift right
	00000	00010	0	3	arithmetic shift right
	00000	00001	0	2	arithmetic shift right
	01101	00001	0		A = A - M
	00110	10000	1	1	arithmetic shift right
	11001	10000	1		A = A + M
	11100	11000	0	0	arithmetic shift right

Answer is  $1110011000_2$ , which is  $-0001101000_2 = -104_{10}$

5. Similar to the above; the quotient is in  $Q = 00010$  and the remainder is in  $A = 00101$ . But now you must figure out the signs, as the original division was  $\frac{17}{-6}$ . Use the equation:

$$\text{Dividend} = \text{Divisor} \times \text{Quotient} + \text{Remainder}$$

To make this work, the quotient is -2 and the remainder is +5:

$$17 = -6 \times -2 + 5$$